

Parsing Techniques
for
Lexicalized
Context-Free Grammars*

Summary

- Part I: Lexicalized Context-Free Grammars
 - motivations and definition
 - relation with other formalisms
- Part II: standard parsing
 - TD techniques
 - BU techniques
- Part III: novel algorithms
 - BU enhanced
 - TD enhanced

Lexicalized grammars

- each rule specialized for one or more lexical items
- advantages over non-lexicalized formalisms:
 - express **syntactic preferences** that are sensitive to lexical words
 - control **word selection**

Syntactic preferences

- adjuncts

Workers [dumped sacks] into a bin

*Workers dumped [sacks into a bin]

- N-N compound

[hydrogen ion] exchange

*hydrogen [ion exchange]

Word selection

- lexical

Nora convened the meeting

?Nora convened the party

- semantics

Peggy solved two puzzles

?Peggy solved two goats

- world knowledge

Mary shelved some books

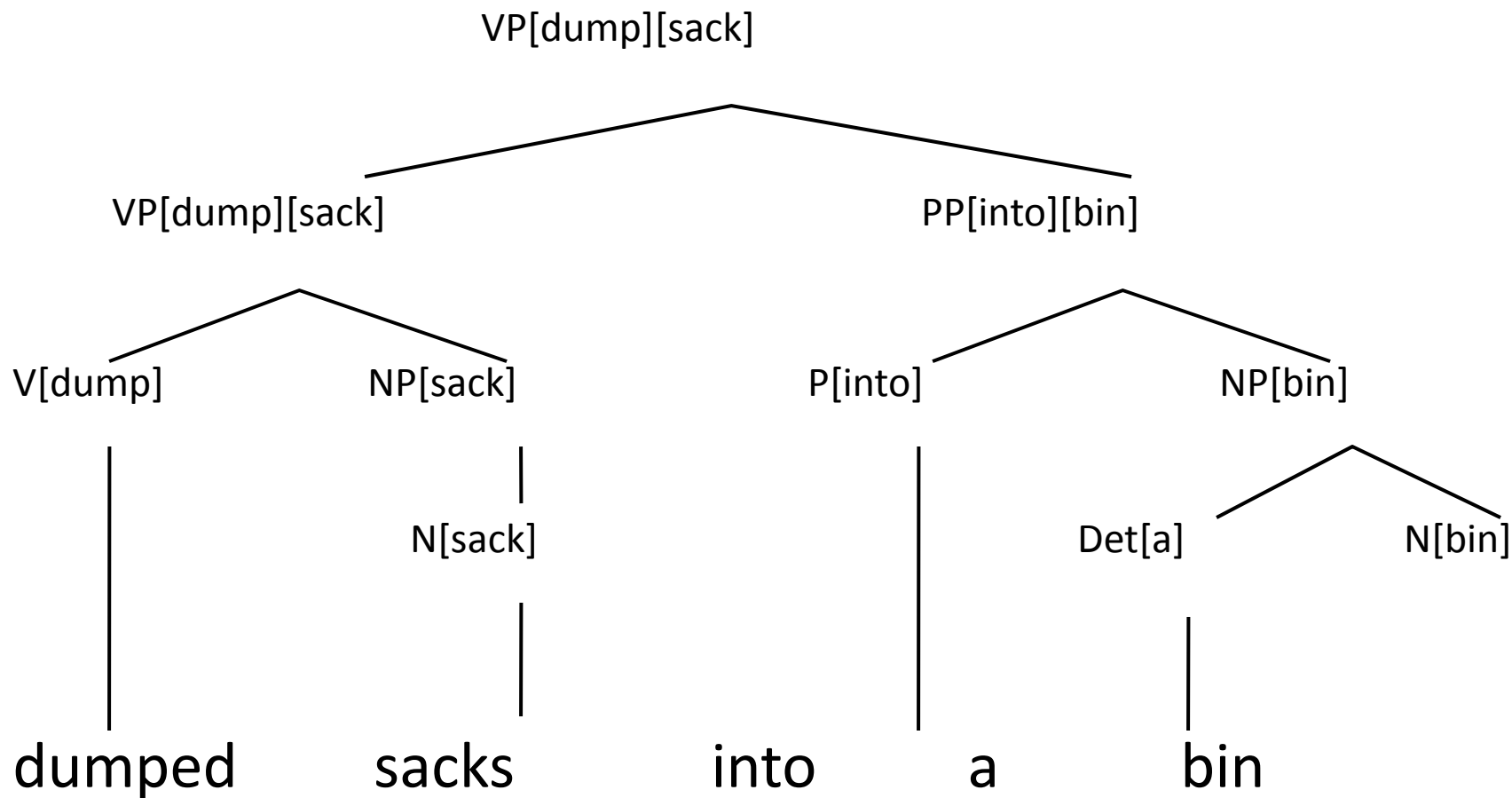
?Mary shelved some cooks

Lexicalized CFG

Motivations :

- study computational properties common to generative formalisms used in state-of-the-art real-world parsers
- develop parsing algorithm that can be directly applied to these formalisms

Lexicalized CFG



Lexicalized CFG

Context-free grammars with :

- **alphabet** V_T :
 - dumped, sacks, into, ...
- **delexicalized nonterminals** V_D :
 - NP, VP, ...
- **nonterminals** V_N :
 - NP[sack], VP[dump][sack], ...

Lexicalized CFG

Delexicalized nonterminals encode :

- word sense
N, V, ...
- grammatical features
number, tense, ...
- structural information
bar level, subcategorization state, ...
- other constraints
distribution, contextual features, ...

Lexicalized CFG

- productions have two forms :
 - V[dump] → dumped
 - VP[dump][sack]
 - VP[dump][sack] PP[into][bin]
- lexical elements in lhs inherited from rhs

Lexicalized CFG

- production is **k-lexical** :
k occurrences of lexical elements in rhs
 - NP[bin] → Det[a] N[bin]
is 2-lexical
 - VP[dump][sack]
→ VP[dump][sack] PP[into][bin]
is 4-lexical

LCFG at work

- 2-lexical CFG
 - Alshawi 1996 : Head Automata
 - Eisner 1996 : Dependency Grammars
 - Charniak 1997 : CFG
 - Collins 1997 : generative model

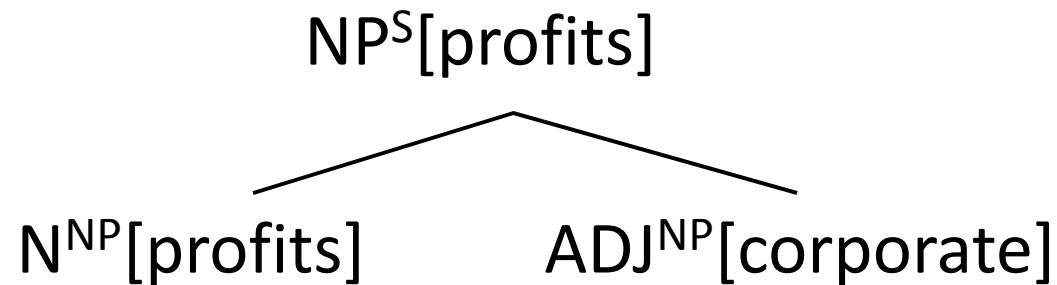
LCFG at work

Probabilistic LCFG G is strongly equivalent to probabilistic grammar G' iff

- 1-2-1 mapping between derivations
- each direction is a homomorphism
- derivation probabilities are preserved

LCFG at work

From Charniak 1997 to 2-lex CFG :



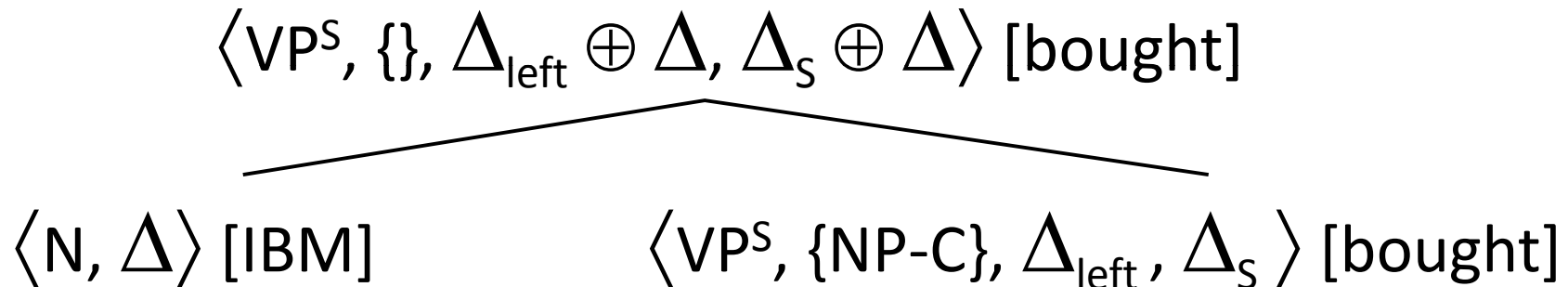
Pr_1 (corporate | ADJ, NP, profits)

Pr_1 (profits | N, NP, profits)

Pr_2 (NP \rightarrow ADJ N | NP, S, profits)

LCFG at work

From Collins 1997 (Model #2) to 2-lex CFG :

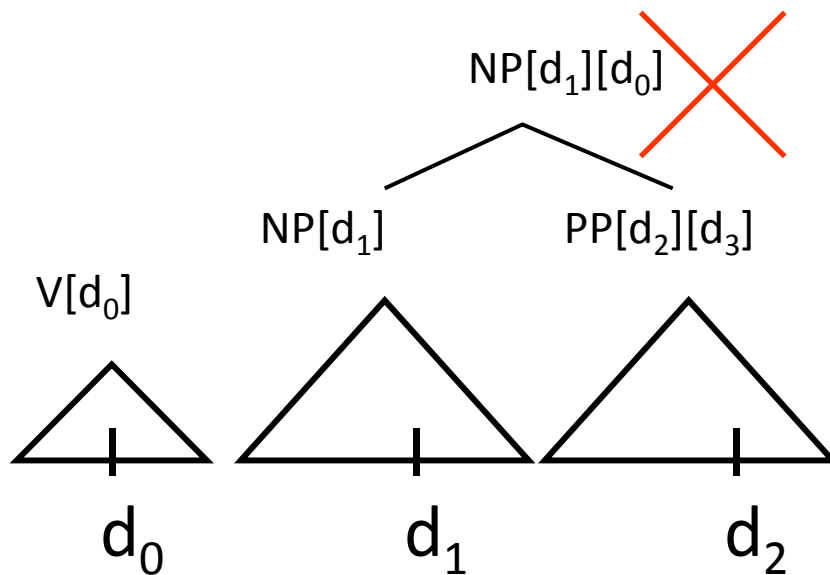


$\text{Pr}_{\text{left}} (\text{NP}, \text{IBM} \mid \text{VP}, \text{S}, \text{bought}, \Delta_{\text{left}}, \{\text{NP-C}\})$

LCFG at work

Major Limitation :

Cannot capture relations involving lexical items **outside** actual constituent (cfr. history based models)



cannot look at d_0
when computing
PP attachment

LCFG at work

- lexicalized context-free parsers that are **not** LCFG :
 - Magerman 1995 : Shift-Reduce+
 - Ratnaparkhi 1997 : Shift-Reduce+
 - Chelba & Jelinek 1998 : Shift-Reduce+
 - Hermjakob & Mooney 1997 : LR

Standard Parsing

- standard parsing algorithms (CKY, Earley, LC, ...) run on LCFG in time $O(|G| \times |w|^3)$
- for 2-lex CFG (simplest case) $|G|$ grows with $|V_D|^3 \times |V_T|^2$!!

Goal :

Get rid of $|V_T|$ factors

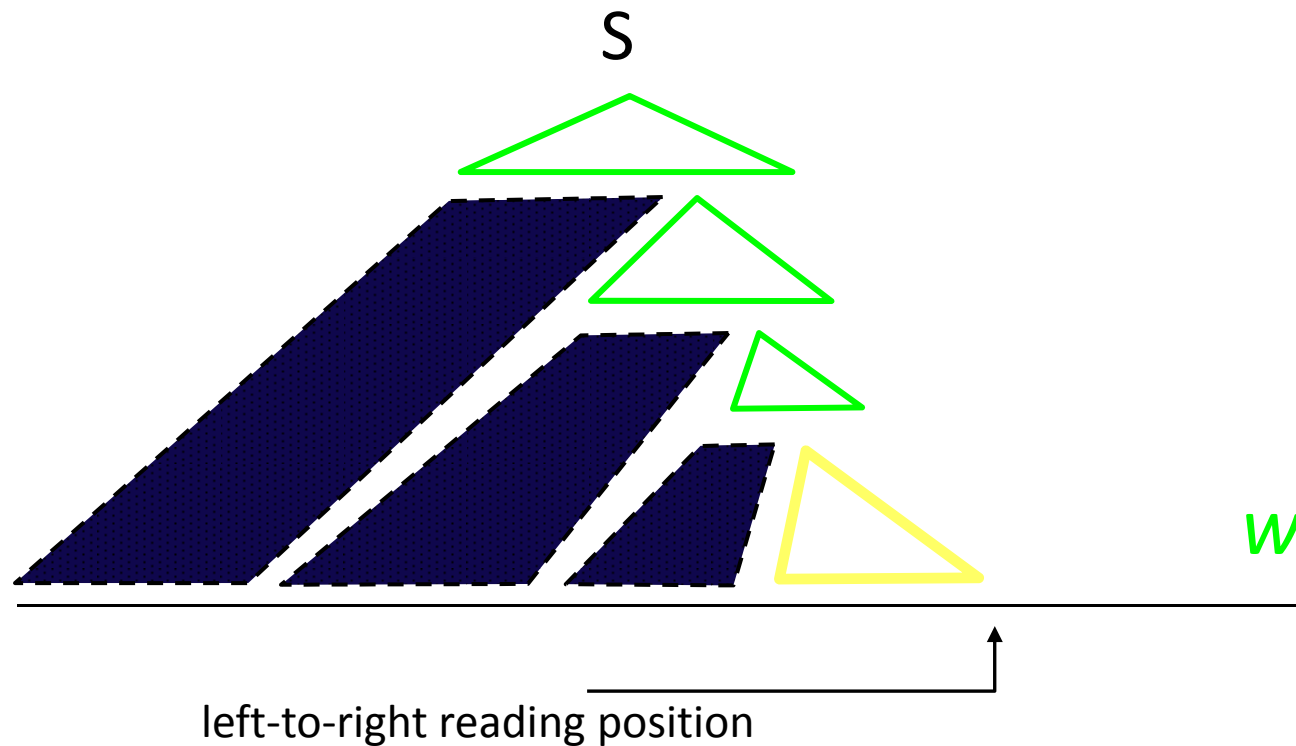
Standard Parsing: TD

Result (to be refined) :

Algorithms satisfying the **correct-prefix property** are “unlikely” to run on LCFG in time independent of V_T

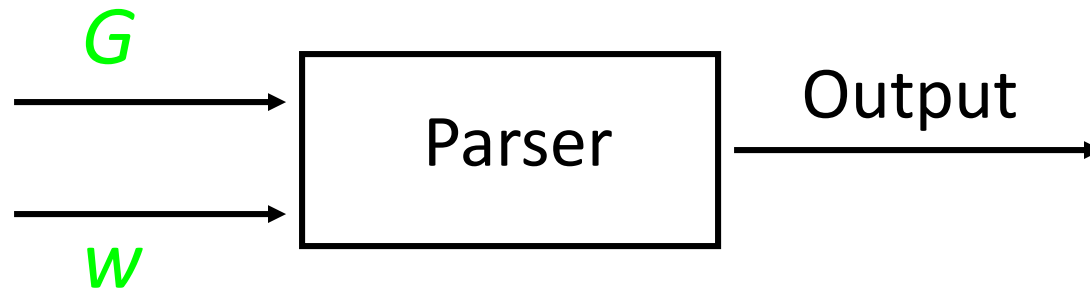
Correct-prefix property

Earley, Left-Corner, GLR, ... :



On-line parsing

No grammar precompilation (Earley) :



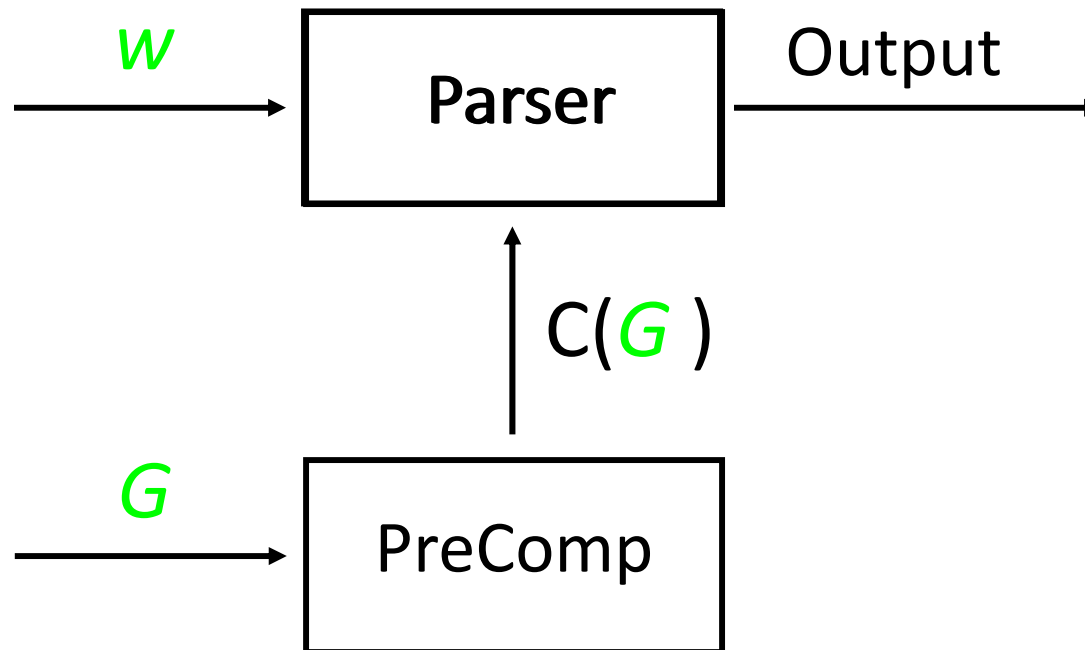
Standard Parsing: TD

Result :

On-line parsers with correct-prefix property cannot run in time $O(f(|V_D|, |w|))$, for any function f

Off-line parsing

Grammar is precompiled (Left-Corner, LR) :



Standard Parsing: TD

Fact :

We can simulate a nondeterministic
FA M on w in time $O(|M| \times |w|)$

Conjecture :

Fix a polynomial p .

We cannot simulate M on w in time
 $p(|w|)$ unless we spend exponential
time in precompiling M

Standard Parsing: TD

Assume our conjecture holds true

Result :

Off-line parsers with correct-prefix property cannot run in time $O(p(|V_D|, |w|))$, for any polynomial p , unless we spend exponential time in precompiling G

Standard Parsing: BU

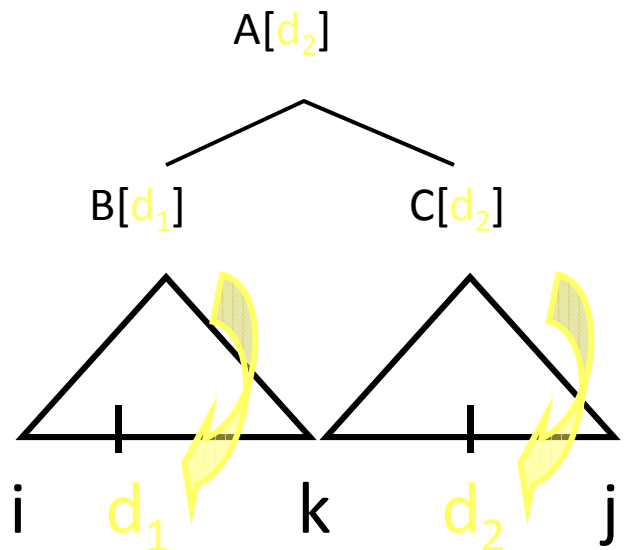
Common practice in lexicalized grammar parsing :

- select productions that are lexically grounded in w
- parse BU with selected subset of G

Problem :

Algorithm removes $|V_T|$ factors but introduces new $|w|$ factors !!

Standard Parsing: BU

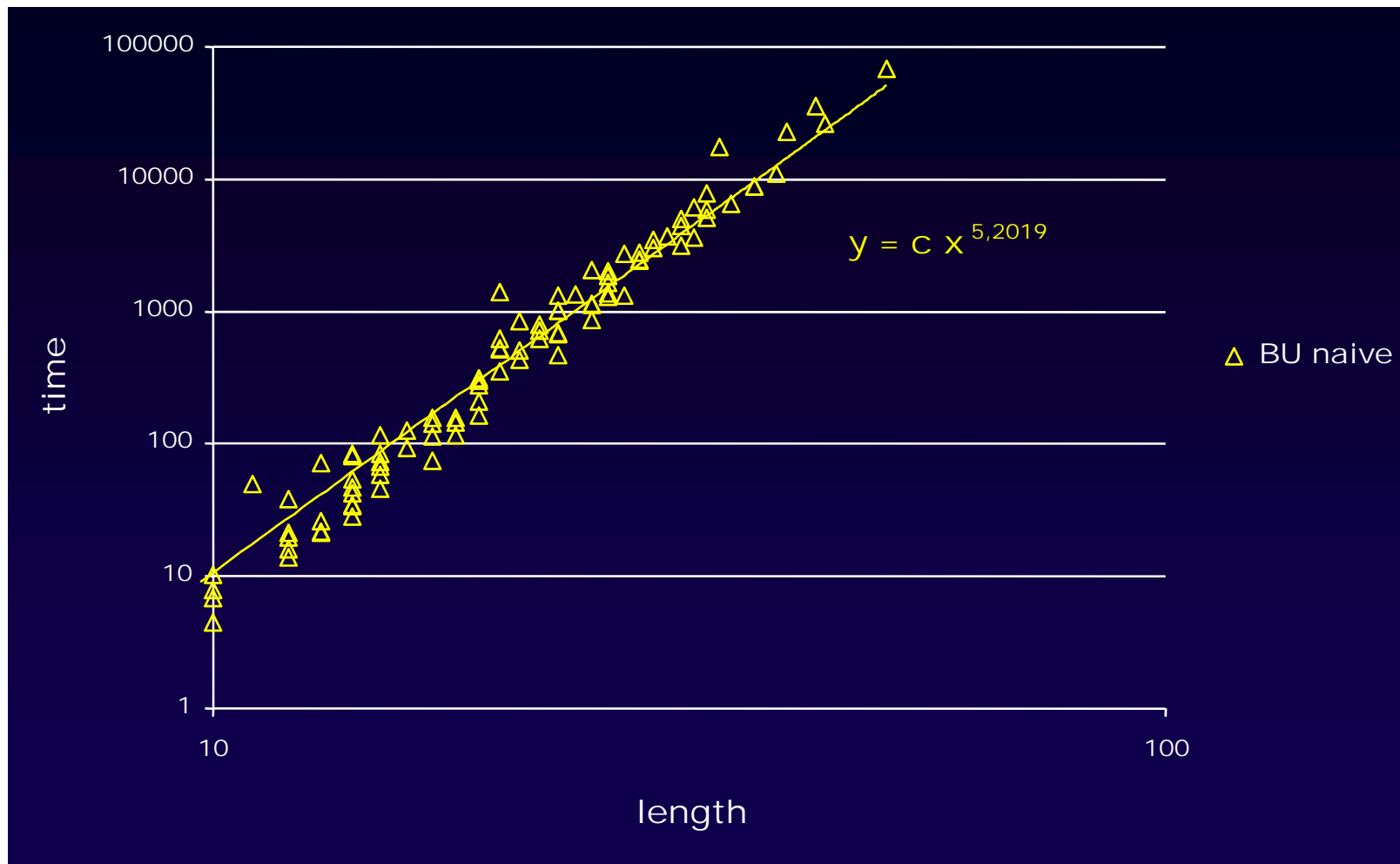


Time charged :

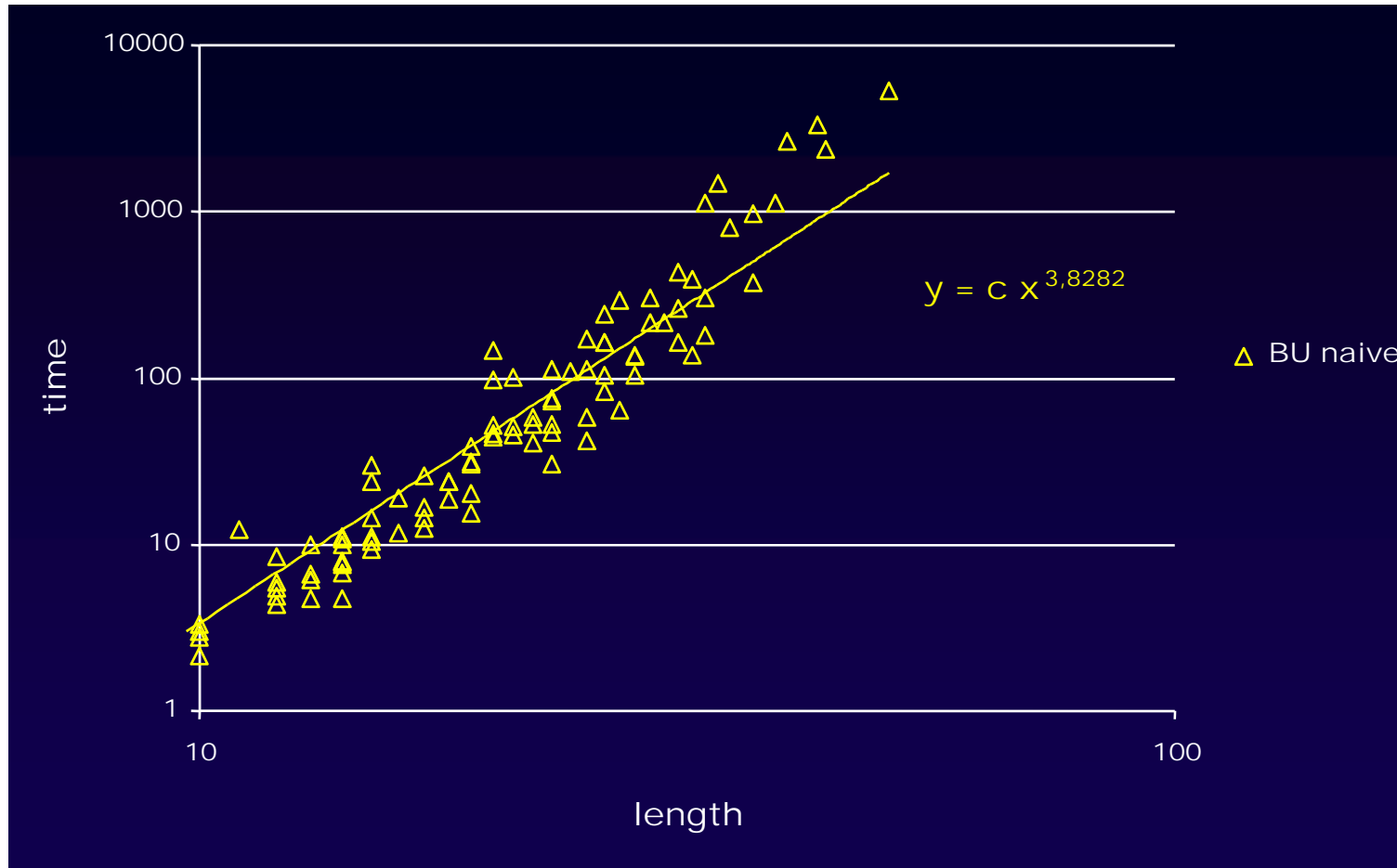
- $i, k, j \Rightarrow |w|^3$
- $A, B, C \Rightarrow |V_D|^3$
- $d_1, d_2 \Rightarrow |w|^2$

Running time is $O(|V_D|^3 \times |w|^5)$!!

Standard BU : Exhaustive



Standard BU : Pruning



novel algorithms

BU enhanced

Result :

Parsing with 2-lex CFG in time

$$O(|V_D|^3 \times |w|^4)$$

Remark :

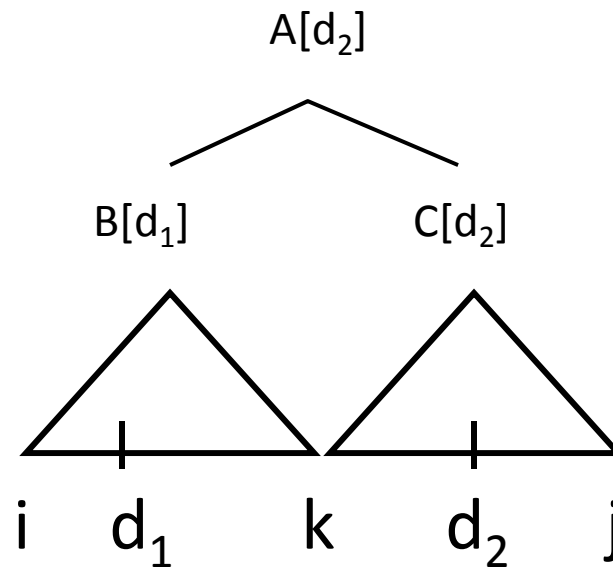
Result transfers to models in Alshawi 1996,
Eisner 1996, Charniak 1997, Collins 1997

Remark :

Technique extends to improve parsing of
Lexicalized-Tree Adjoining Grammars

Algorithm #1

Basic step in naive BU :

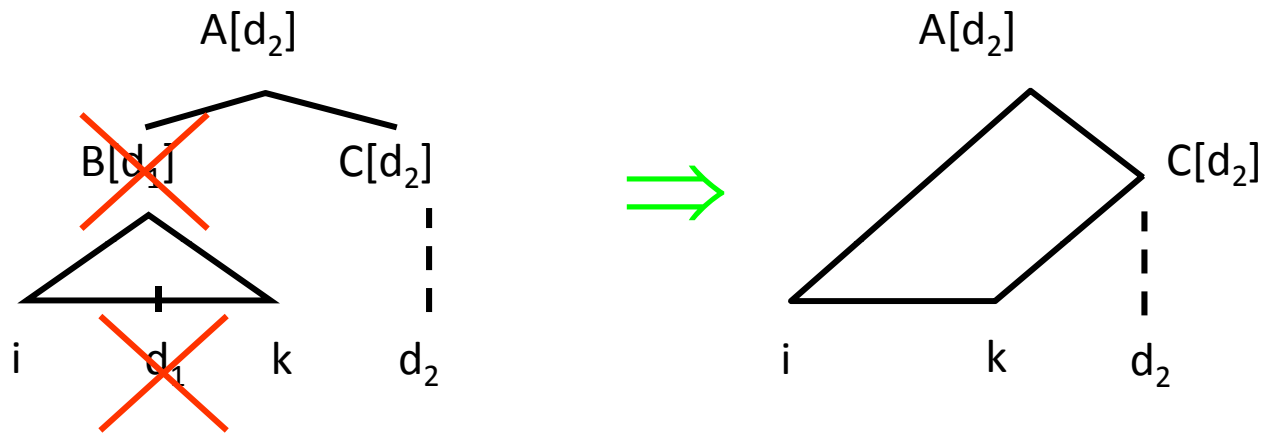


Idea :

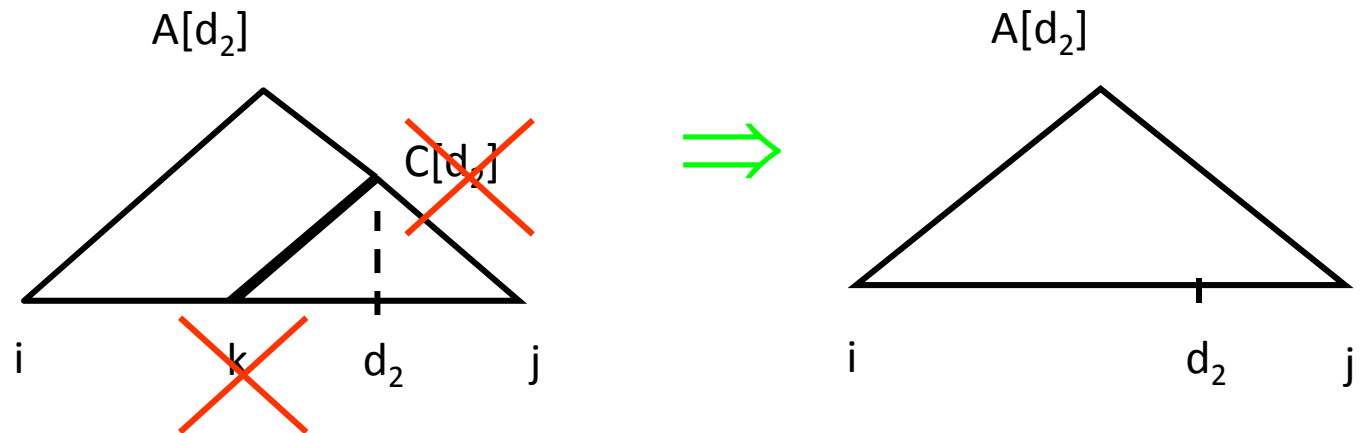
Indices d_1 and j can be processed independently

Algorithm #1

- Step 1



- Step 2



BU enhanced

Upper bound provided by Algorithm #1 :

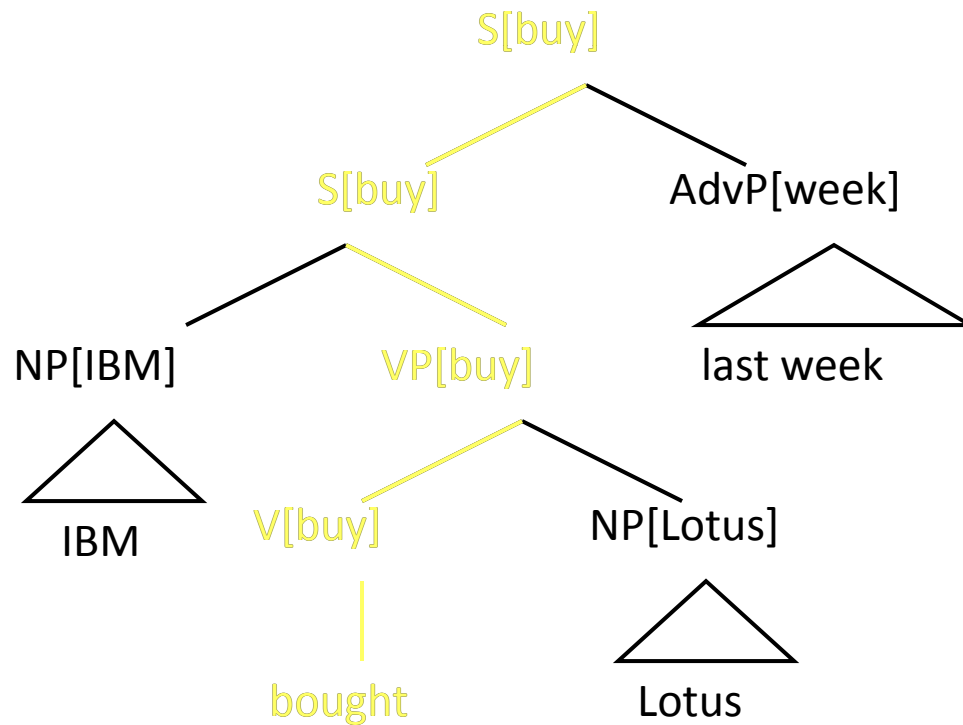
$$O(|w|^4)$$

Goal :

Can we go down to $O(|w|^3)$?

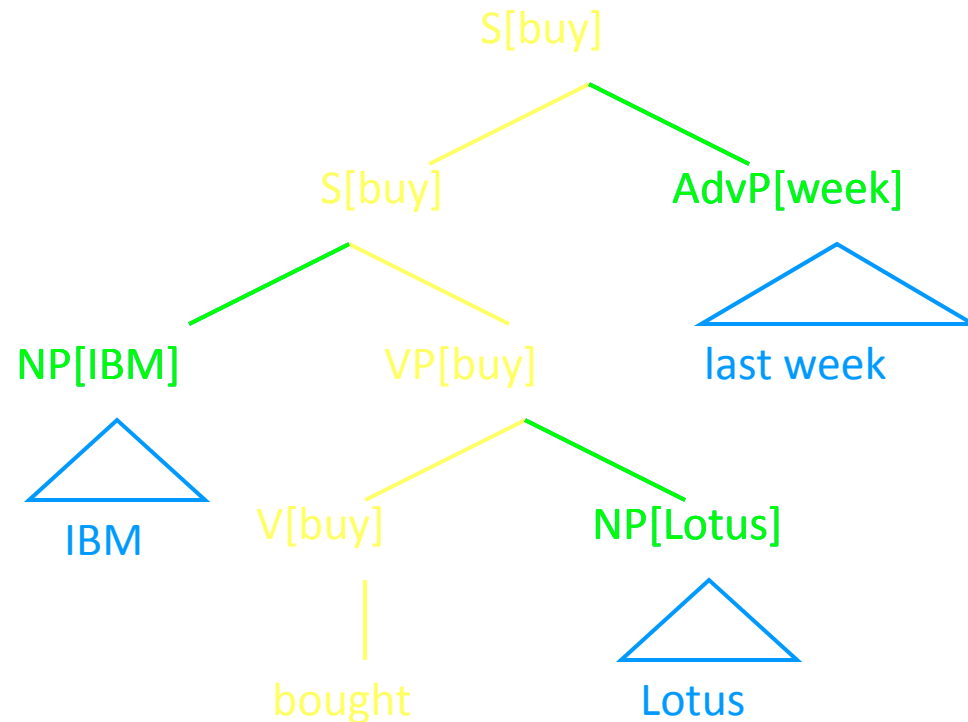
Spine

The **spine** of a parse tree is the path from the root to the root's head



Spine projection

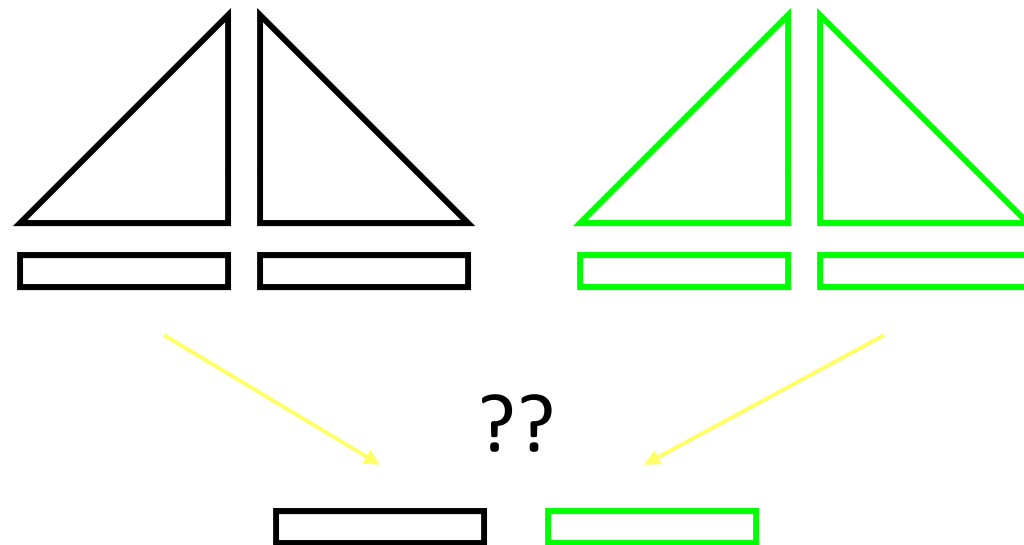
The **spine projection** is the yield of the sub-tree composed by the spine and all its sibling nodes



NP[IBM] bought NP[Lotus] AdvP[week]

Split Grammars

Split spine projections at head :



Problem :

how much information do we need to store in order to construct new grammatical spine projections from splits ?

Split Grammars

Fact :

Set of spine projections is a linear context-free language

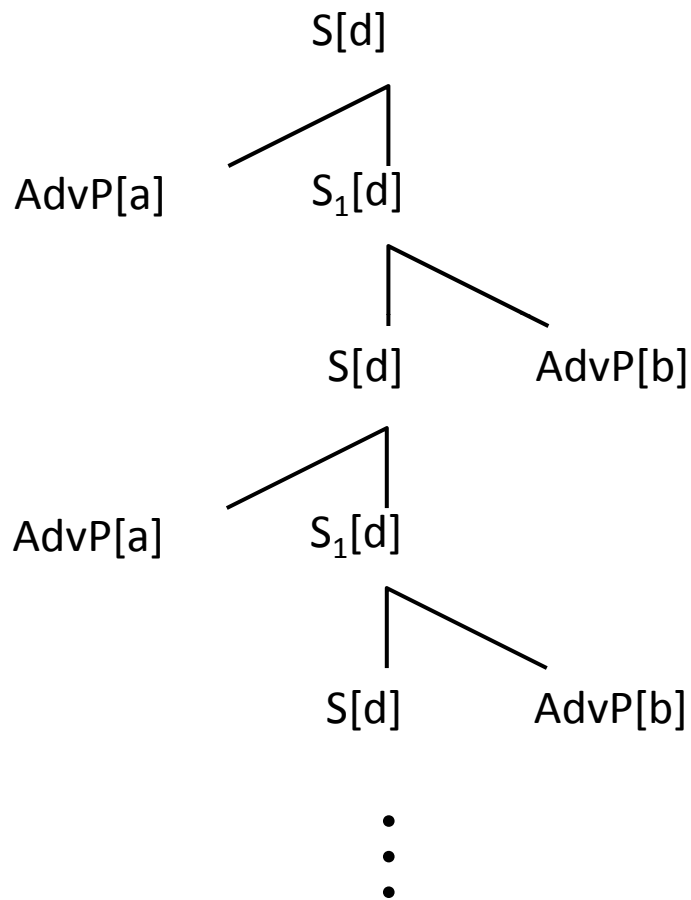
Definition :

2-lex CFG is **split** if set of spine projections is a regular language

Remark :

For split grammars, we can recombine splits using finite information

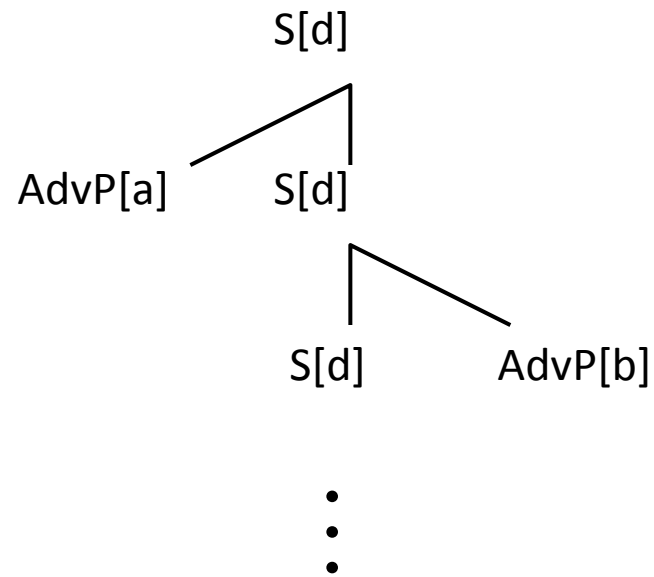
Split Grammars



Non-split grammar :

- unbounded # of dependencies between left and right dependents of head
- linguistically unattested and unlikely

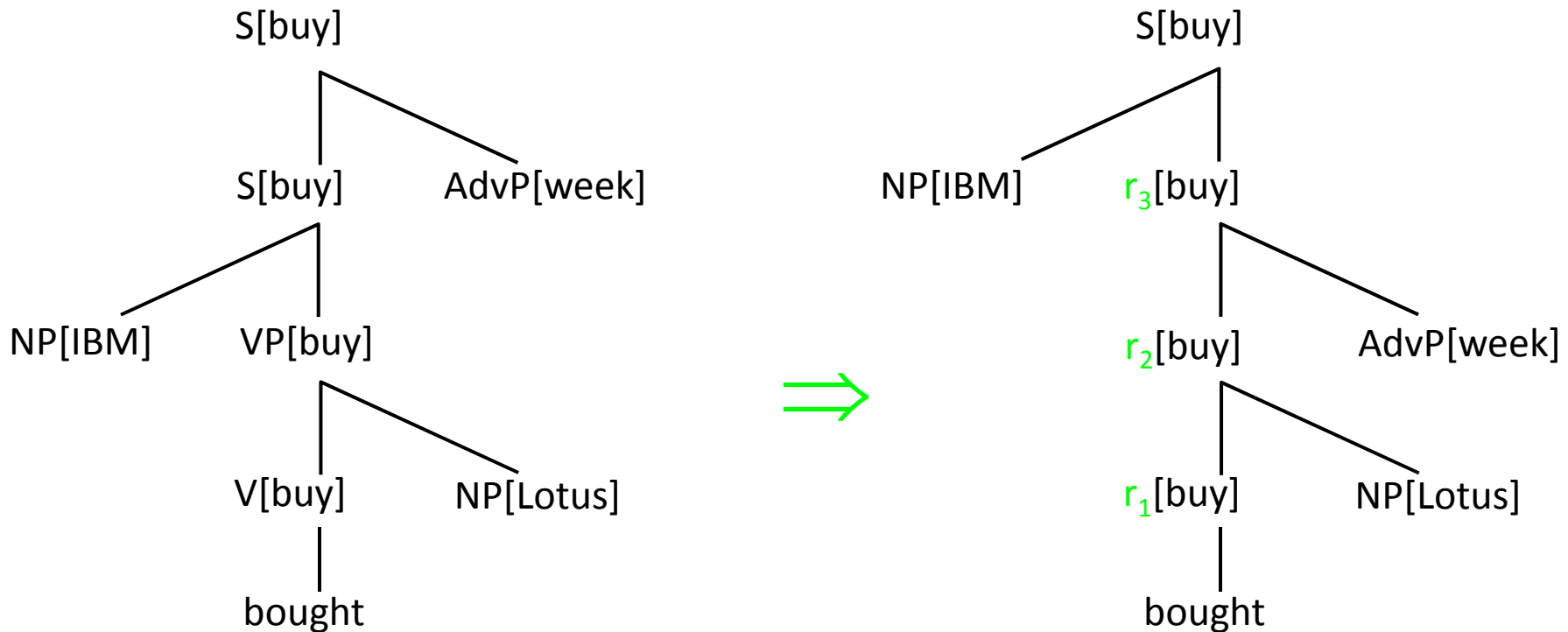
Split Grammars



Split grammar :
finite # of
dependencies
between left and
right dependents
of lexical head

Split Grammars

Precompile grammar such that splits are derived separately



$r_3[\text{buy}]$ is a split symbol

Split Grammars

- t : max # of states per spine automaton
- g : max # of split symbols per spine automaton ($g < t$)
- m : # of delexicalized nonterminals there are maximal projections

BU enhanced

Result :

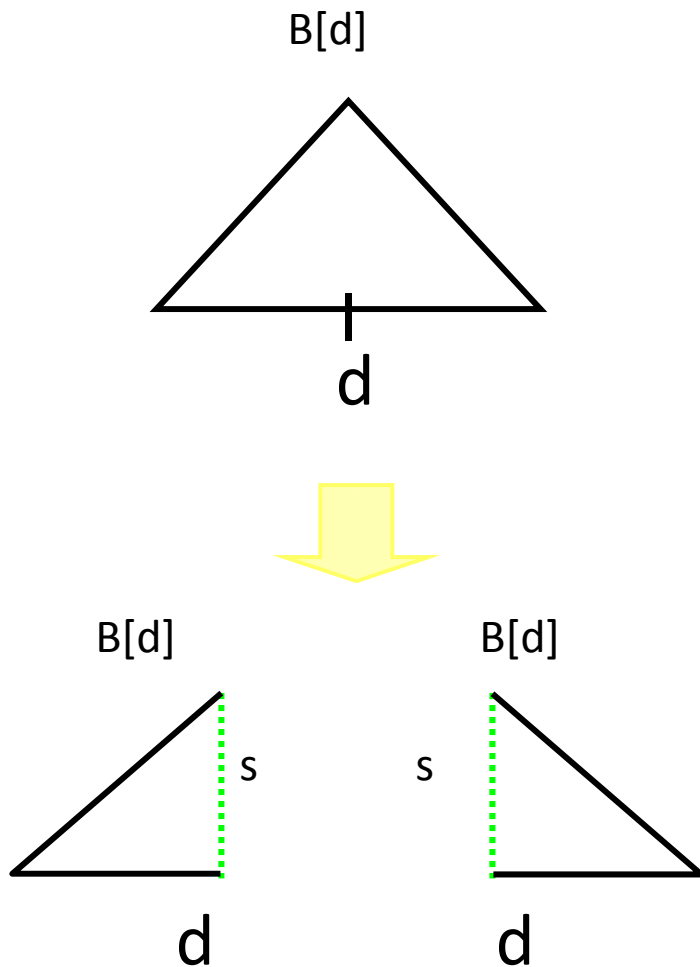
Parsing with split 2-lexical CFG in time

$$O(t^2 g^2 m^2 |w|^3)$$

Remark:

Models in Alshawi 1996, Charniak 1997
and Collins 1997 are not split

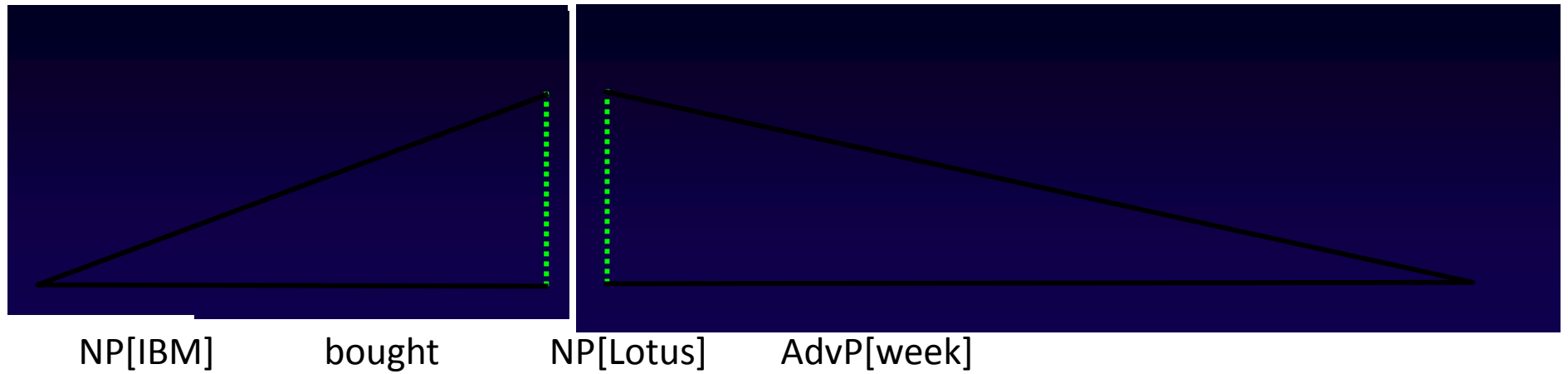
Algorithm #2



Idea :

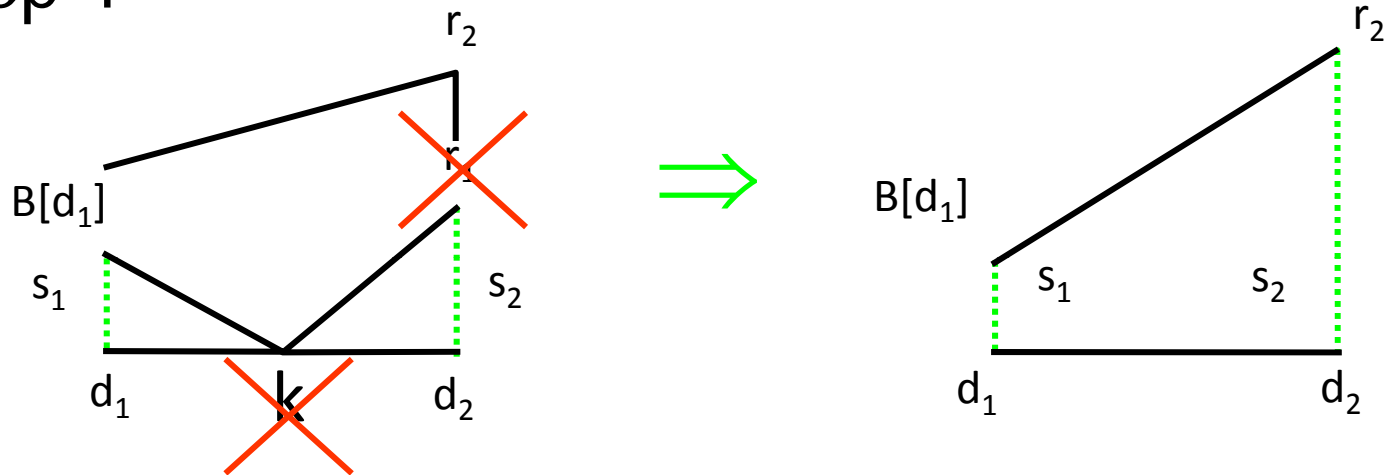
- recognize left and right splits separately
- collect head dependents one split at a time

Algorithm #2

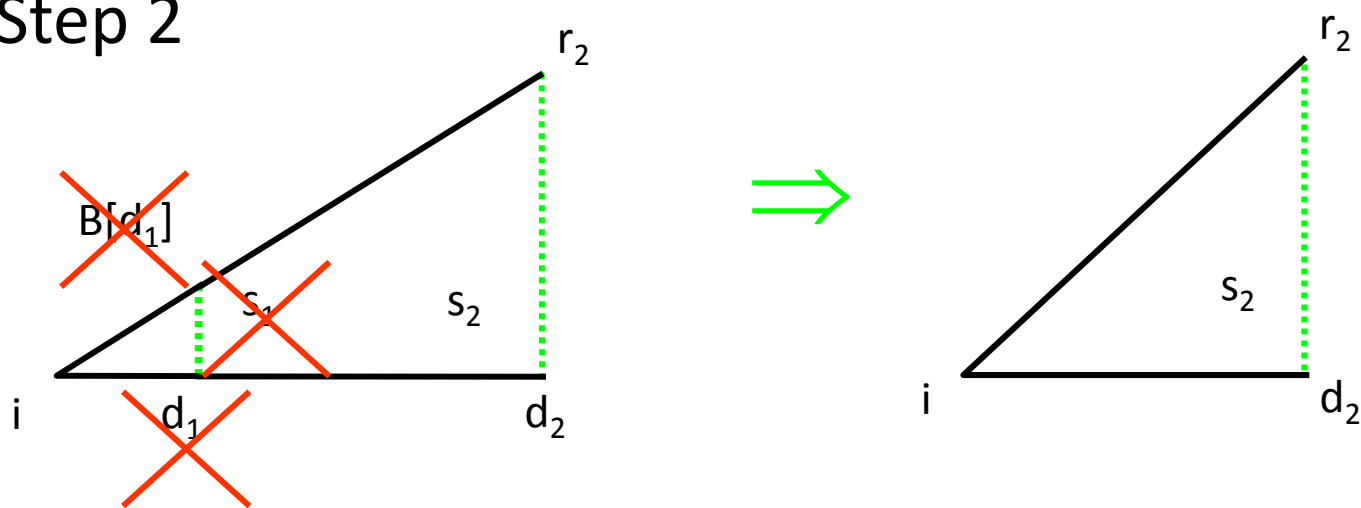


Algorithm #2

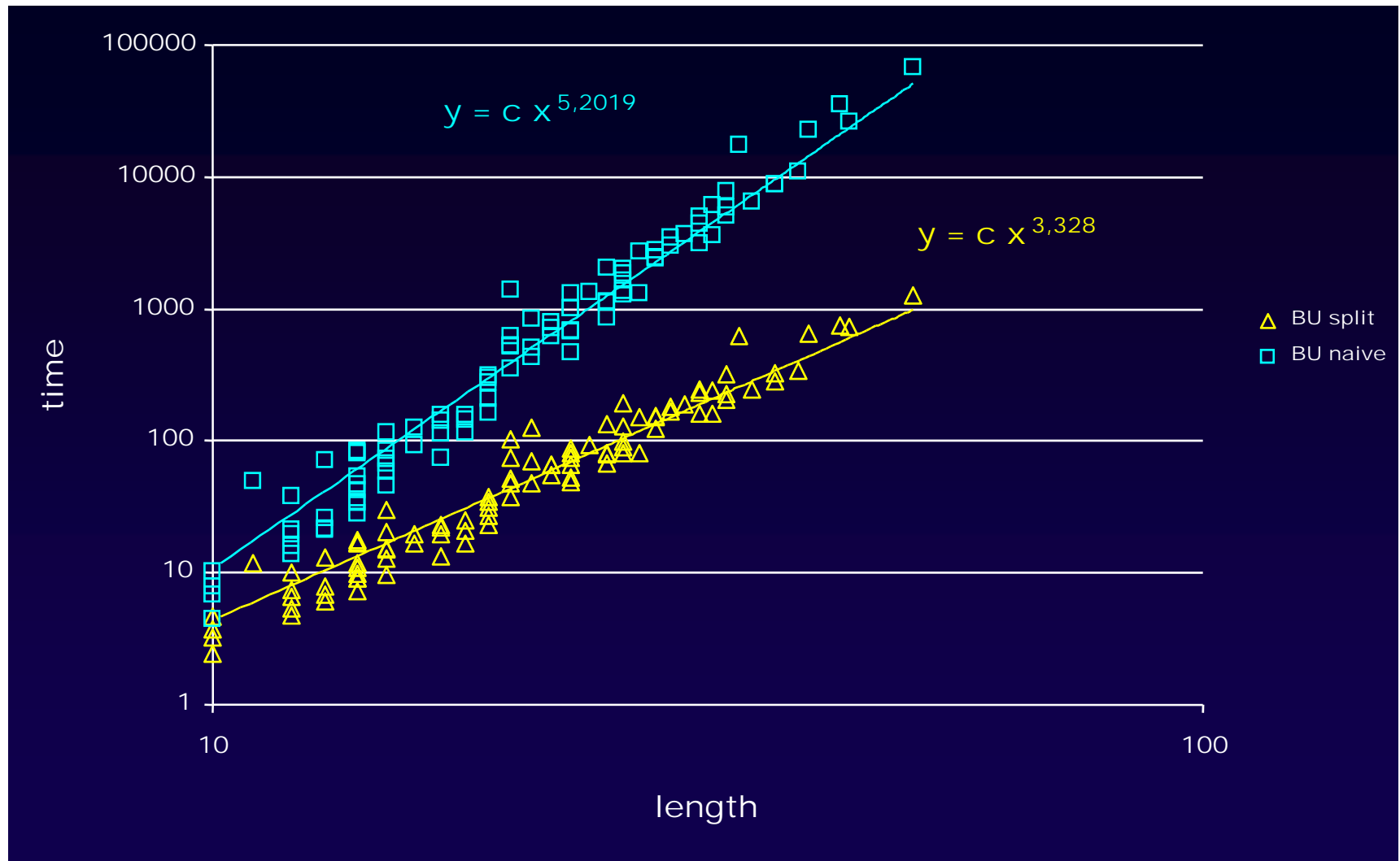
- Step 1



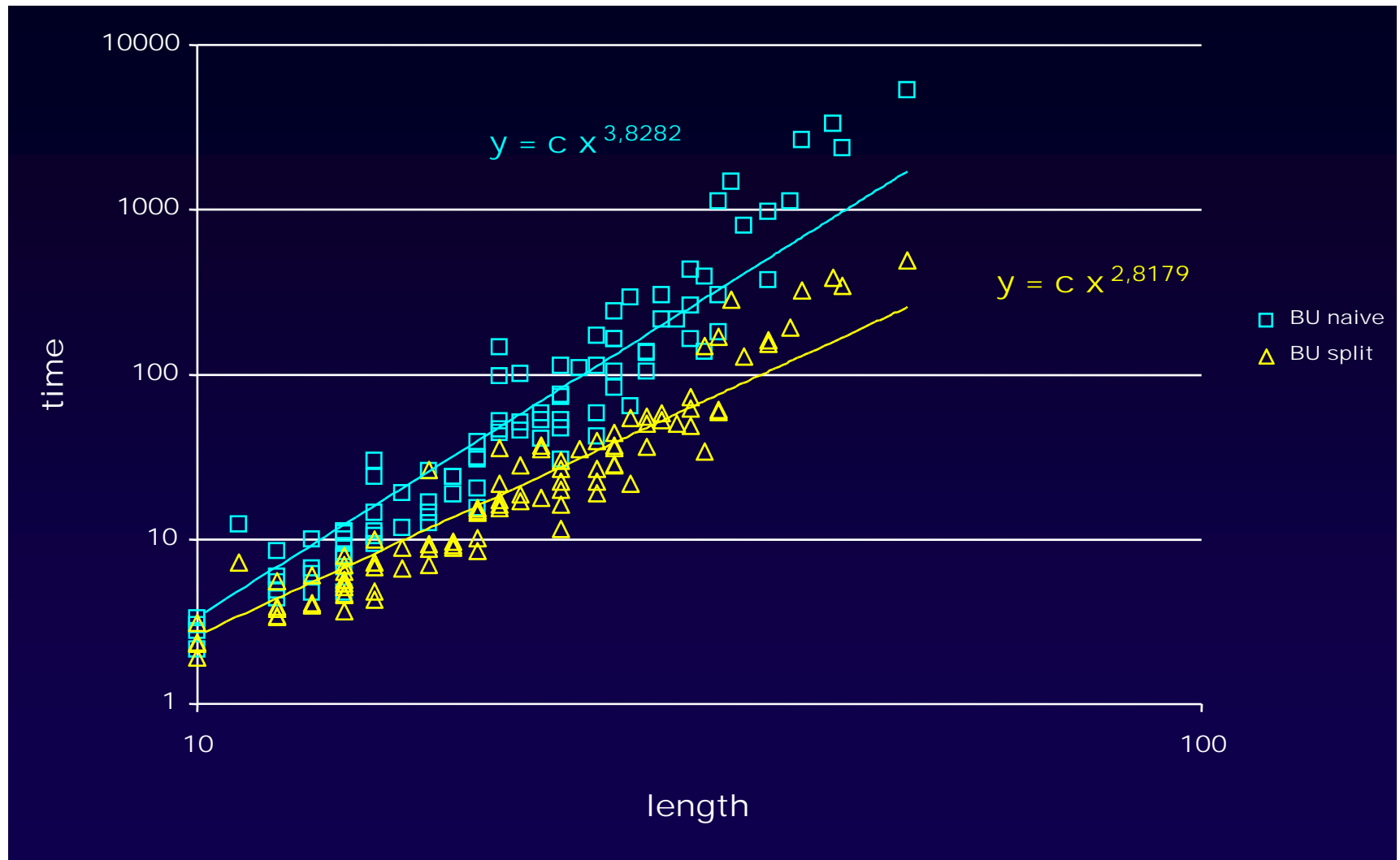
- Step 2



Algorithm #2 : Exhaustive



Algorithm #2 : Pruning



TD enhanced

Goal :

Introduce TD prediction for 2-lexical CFG parsing, without $|V_T|$ factors

Remark :

Must relax left-to-right parsing
(because of previous results)

TD enhanced

Result :

TD parsing with 2-lex CFG in time

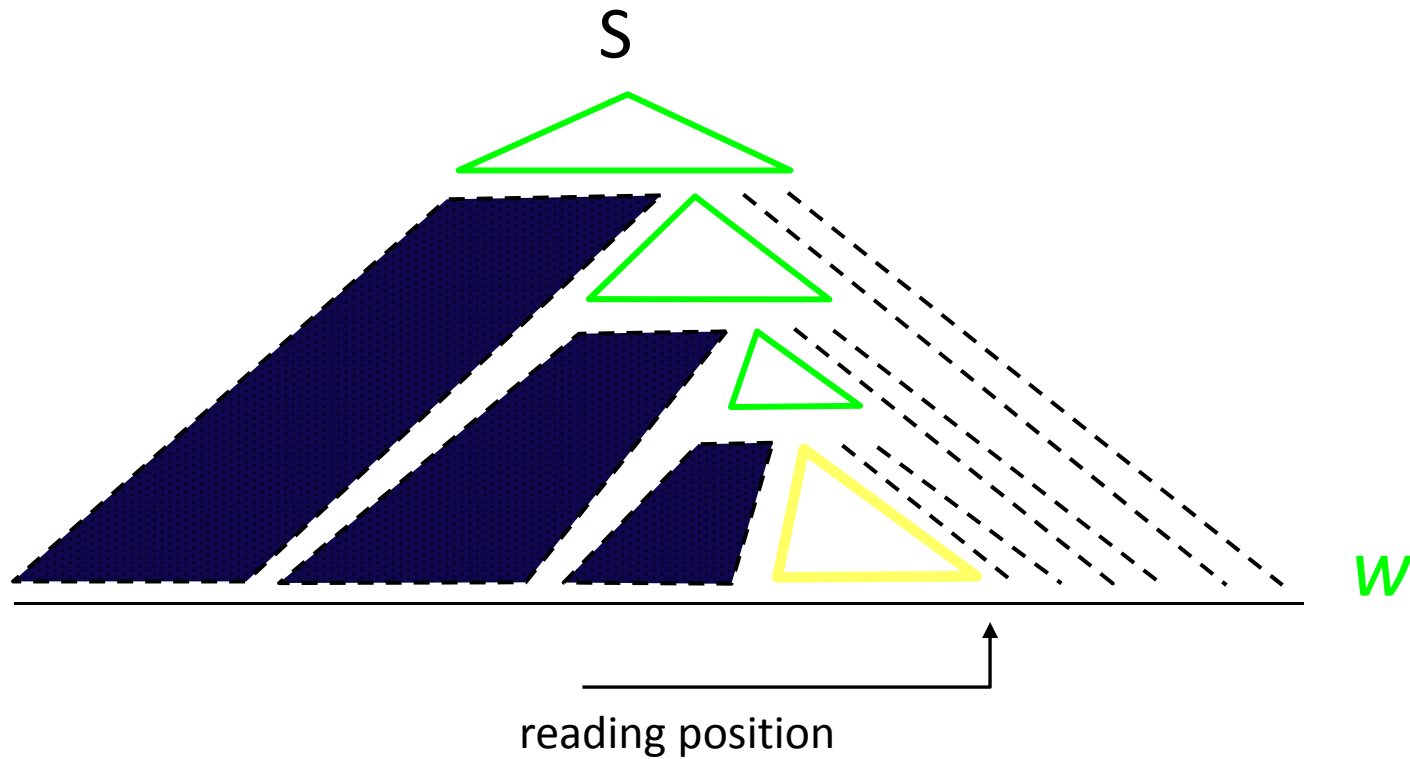
$$O(|V_D|^3 \times |w|^4)$$

Open :

$O(|w|^3)$ extension to split grammars

TD enhanced

Strongest version of correct-prefix property :

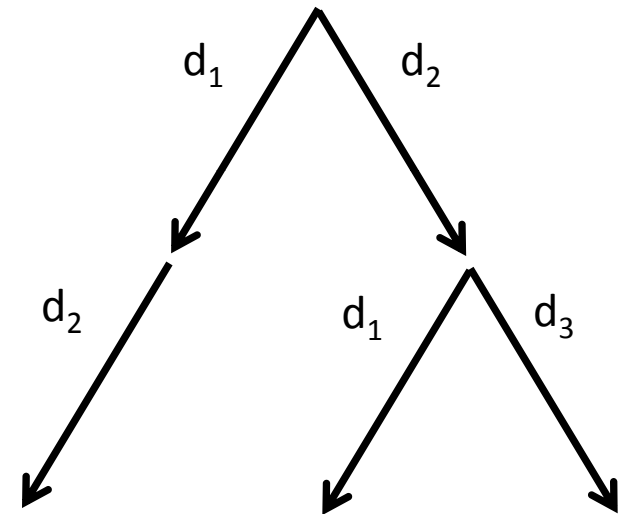


Data Structures

Prods with lhs $A[d]$:

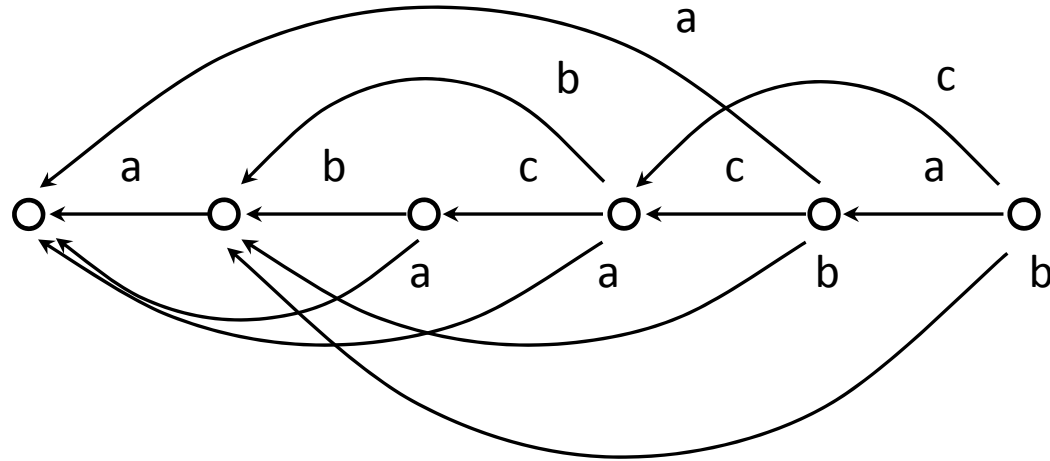
- $A[d] \rightarrow X_1[d_1] X_2[d_2]$
- $A[d] \rightarrow Y_1[d_3] Y_2[d_2]$
- $A[d] \rightarrow Z_1[d_2] Z_2[d_1]$

Trie for $A[d]$:

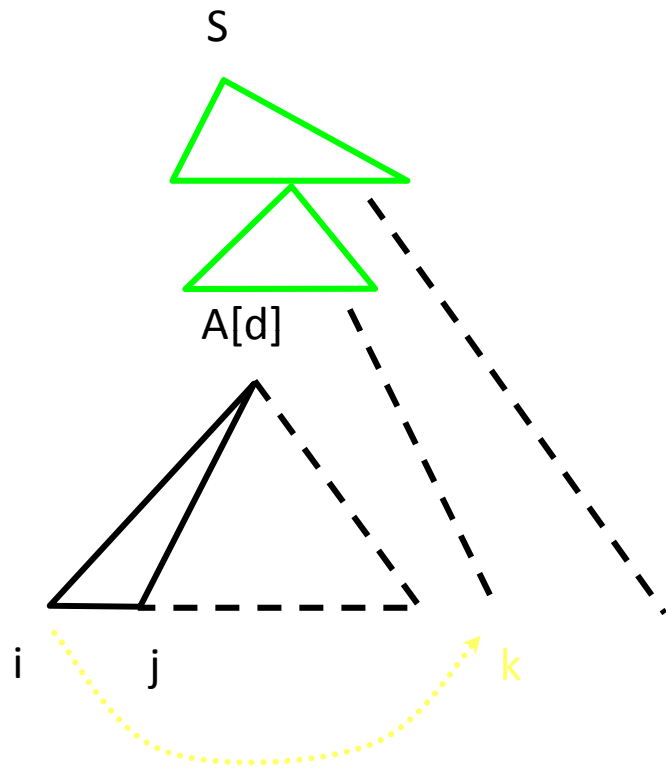


Data Structures

Rightmost subsequence recognition
by precompiling input w into a
deterministic FA :



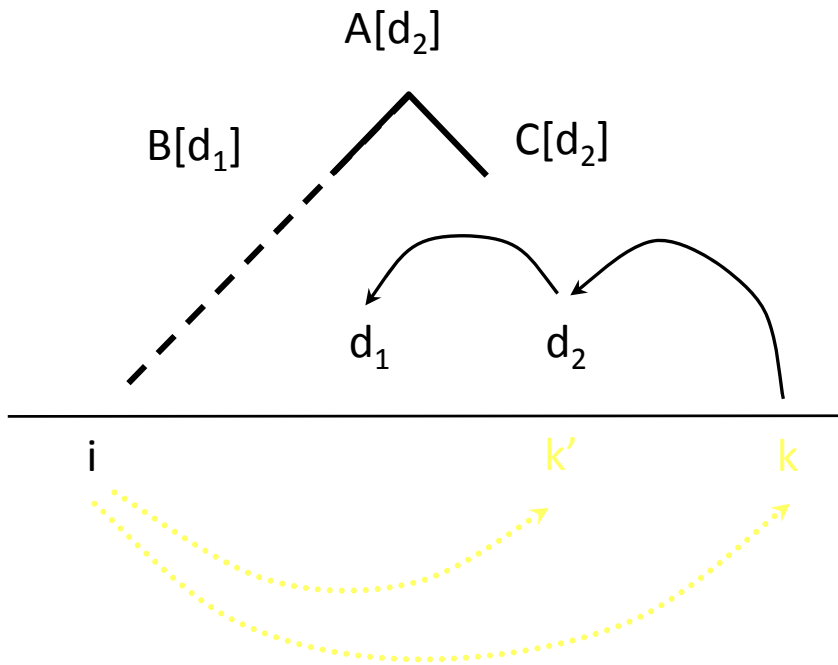
Algorithm #3



Item representation :

- i, j indicate extension of $A[d]$ partial analysis
- k indicates rightmost possible position for completion of $A[d]$ analysis

Algorithm #3 : Prediction



- Step 1 :
find rightmost
subsequence
before k for some
 $A[d_2]$ production
- Step 2 :
make Earley
prediction

Conclusions

- standard parsing techniques are not suitable for processing lexicalized grammars
- novel algorithms have been introduced using enhanced dynamic programming
- work to be done :
extension to history-based models